

Application layer

Goals:

- ❑ Conceptual aspects of network application protocols
 - Client server paradigm
 - Service models
- ❑ Review protocols by examining popular application-level protocols
 - HTTP
 - DNS

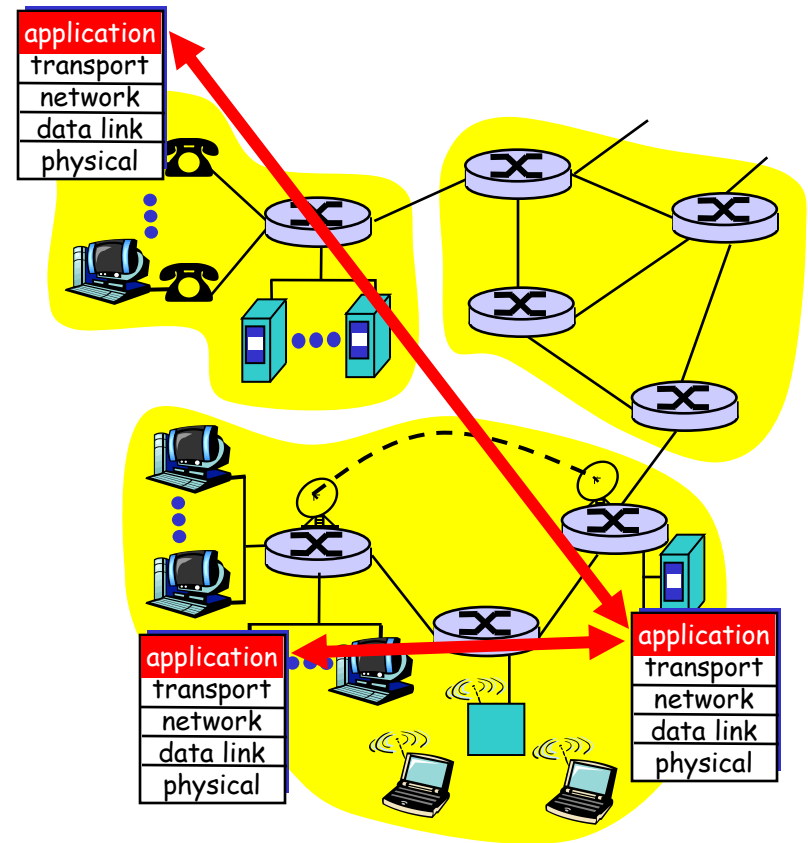
Applications and application-layer protocols

Application: communicating, distributed processes

- Running in network hosts in “user space”
- Exchange messages to implement app
- E.g., email, file transfer, the Web

Application-layer protocols

- One “piece” of an app
- Define messages exchanged by apps and actions taken
- User services provided by lower layer protocols



Client-server paradigm

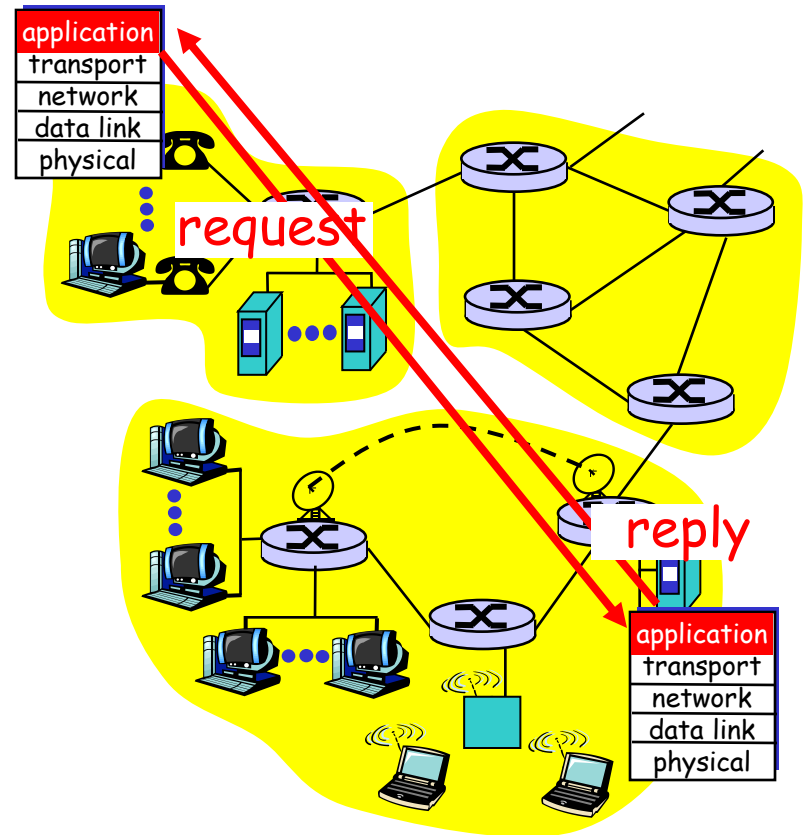
Typical network app has two pieces: *client* and *server*

Client:

- ❑ Initiates contact with server (“speaks first”)
- ❑ Typically requests service from server,
- ❑ E.g., request WWW page, send email

Server:

- ❑ Provides requested service to client
- ❑ E.g., sends requested WWW page, receives/stores received email



Services provided by Internet transport protocols

TCP service:

- ❑ *Connection-oriented*: setup required between client, server
- ❑ *Reliable transport* between sending and receiving process
- ❑ *Flow control*: sender won't overwhelm receiver
- ❑ *Congestion control*: throttle sender when network overloaded
- ❑ *Does not providing*: timing, minimum bandwidth guarantees

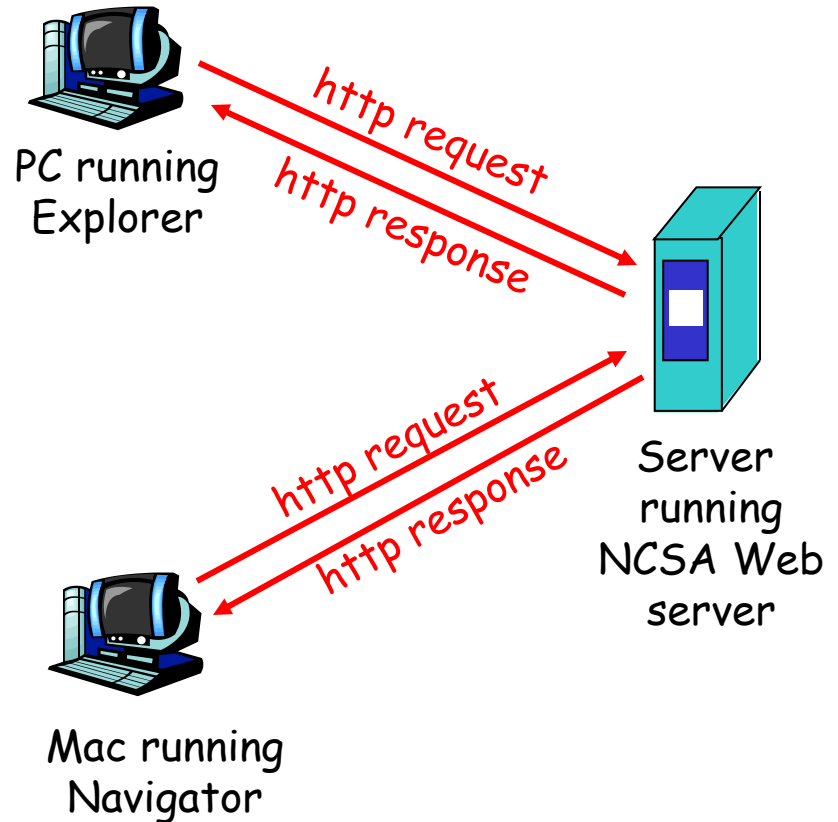
UDP service:

- ❑ Unreliable data transfer between sending and receiving process
- ❑ Does not provide: connection setup, reliability, flow control, congestion control, timing, or bandwidth guarantee

WWW: The HTTP protocol

HTTP: hypertext transfer protocol

- ❑ WWW's application layer protocol
- ❑ Client/server model
 - *Client*: browser that requests, receives, "displays" WWW objects
 - *Server*: WWW server sends objects in response to requests



HTTP - timeline

- ❑ Mar 1990 CERN labs document proposing Web
- ❑ Jan 1992 HTTP/0.9 specification
- ❑ Dec 1992 Proposal to add MIME to HTTP
- ❑ Feb 1993 UDI (Universal Document Identifier) Network
- ❑ Mar 1993 HTTP/1.0 first draft
- ❑ Jun 1993 HTML (1.0 Specification)
- ❑ Oct 1993 URL specification
- ❑ Nov 1993 HTTP/1.0 second draft
- ❑ Mar 1994 URI in WWW
- ❑ May 1996 HTTP/1.0 Informational, RFC 1945
- ❑ Jan 1997 HTTP/1.1 Proposed Standard, RFC 2068
- ❑ Jun 1999 HTTP/1.1 Draft Standard, RFC 2616
- ❑ 2001 HTTP/1.1 Formal Standard
- ❑ ...

The HTTP protocol: More

HTTP: TCP transport service

- ❑ Client initiates TCP connection (creates socket) to server, port 80
- ❑ Server accepts TCP connection from client
- ❑ http messages (application-layer protocol messages) exchanged between browser (http client) and WWW server (http server)
- ❑ TCP connection closed

HTTP is "stateless"

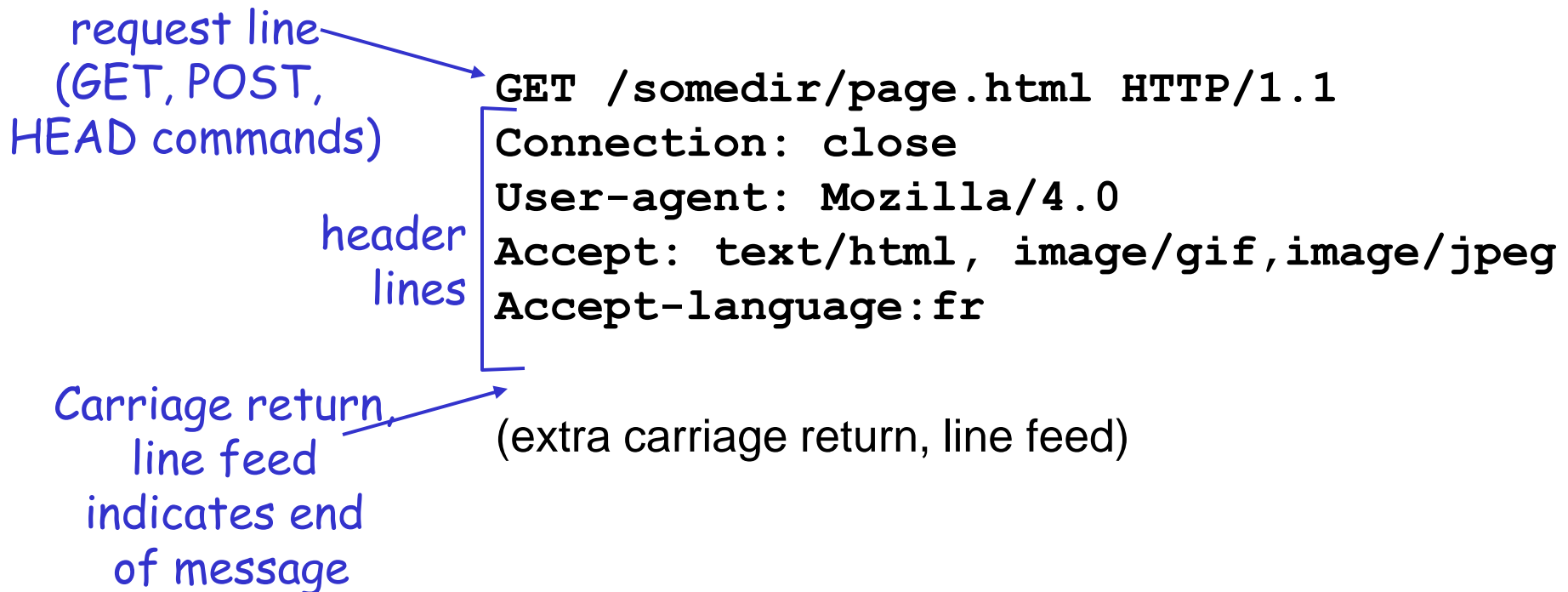
- ❑ Server maintains no information about past client requests

aside
Protocols that maintain "state" are complex!

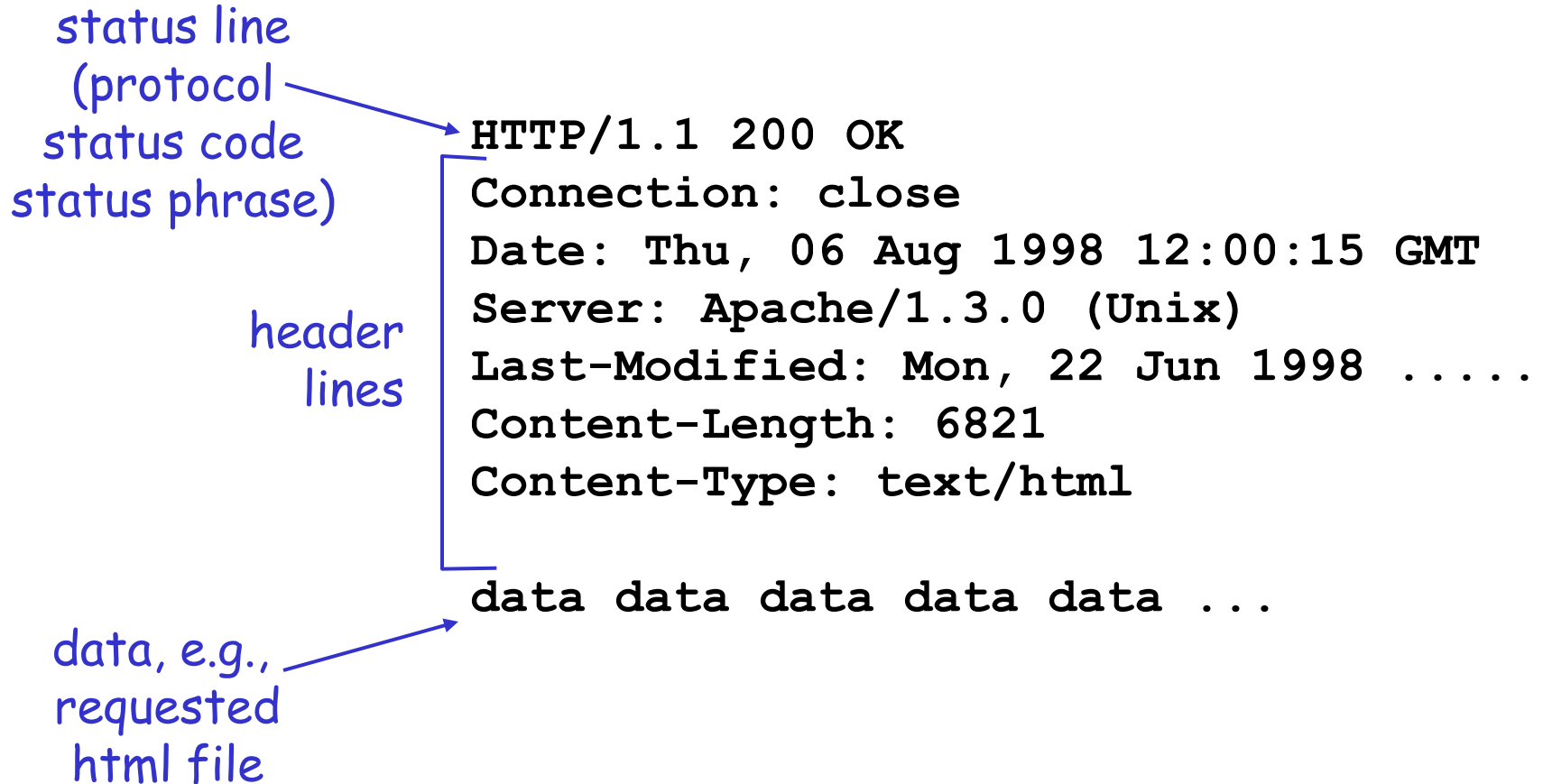
- ❑ Past history (state) must be maintained
- ❑ If server/client crashes, their views of "state" may be inconsistent, must be reconciled

HTTP message format: Request

- ❑ Two types of http messages: *Request, response*
- ❑ **http request message:**
 - ASCII (human-readable format)



HTTP message format: Reply



HTTP reply status codes

In first line in server → client response message.

A few sample codes:

200 OK

- Request succeeded, requested object later in this message

301 Moved Permanently

- Requested object moved, new location specified later in this message (Location:)

400 Bad Request

- Request message not understood by server

404 Not Found

- Requested document not found on this server

505 HTTP Version Not Supported

HTTP request methods

□ Methods

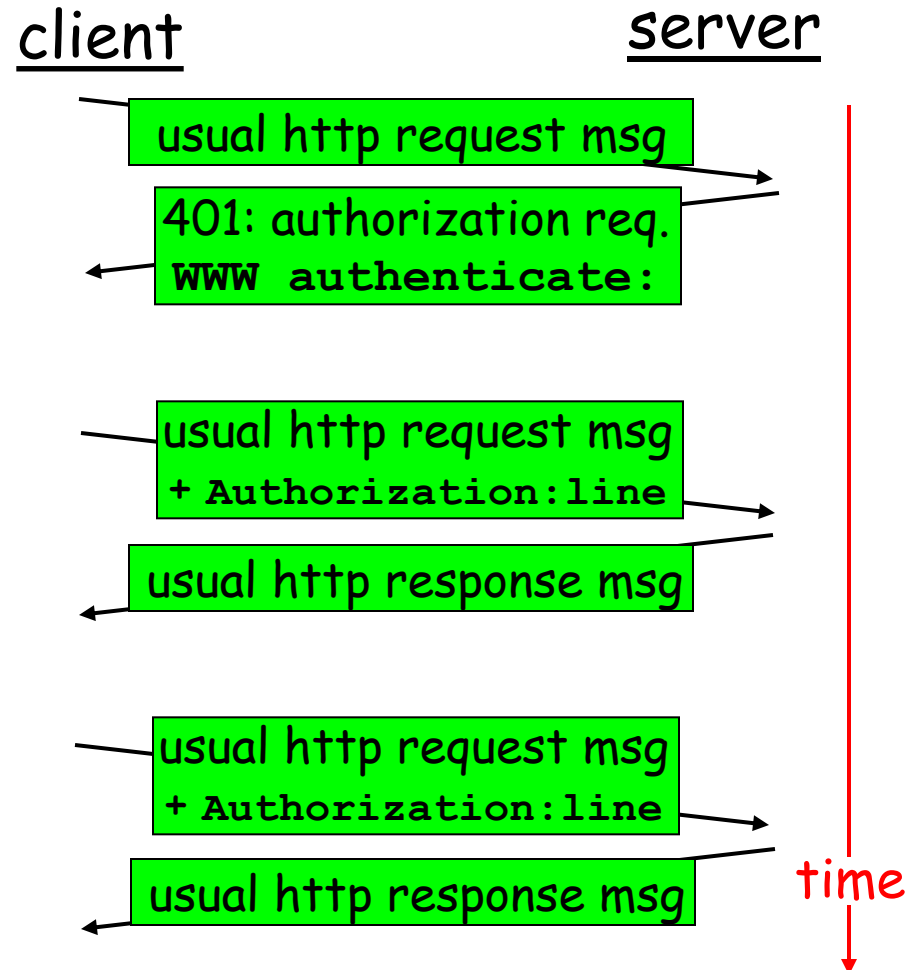
- GET
- HEAD
- POST
- PUT
- Delete

The HTTP protocol: Even more

- ❑ Non-persistent connection:
 - One object in each TCP connection
 - Some browsers create multiple TCP connections *simultaneously* – one per object
- ❑ Persistent connection:
 - Multiple objects transferred within one TCP connection
- ❑ Pipelined persistent connections:
 - Multiple requests issued without waiting for response

User-server interaction: Authentication

- Authentication goal:** Control access to server documents
- ❑ **Stateless:** Client must present authorization in each request
 - ❑ **Authorization:** Typically name, password
 - **Authorization:** header line in request
 - If no authorization, server refuses access, sends `WWW authenticate:` header line in response

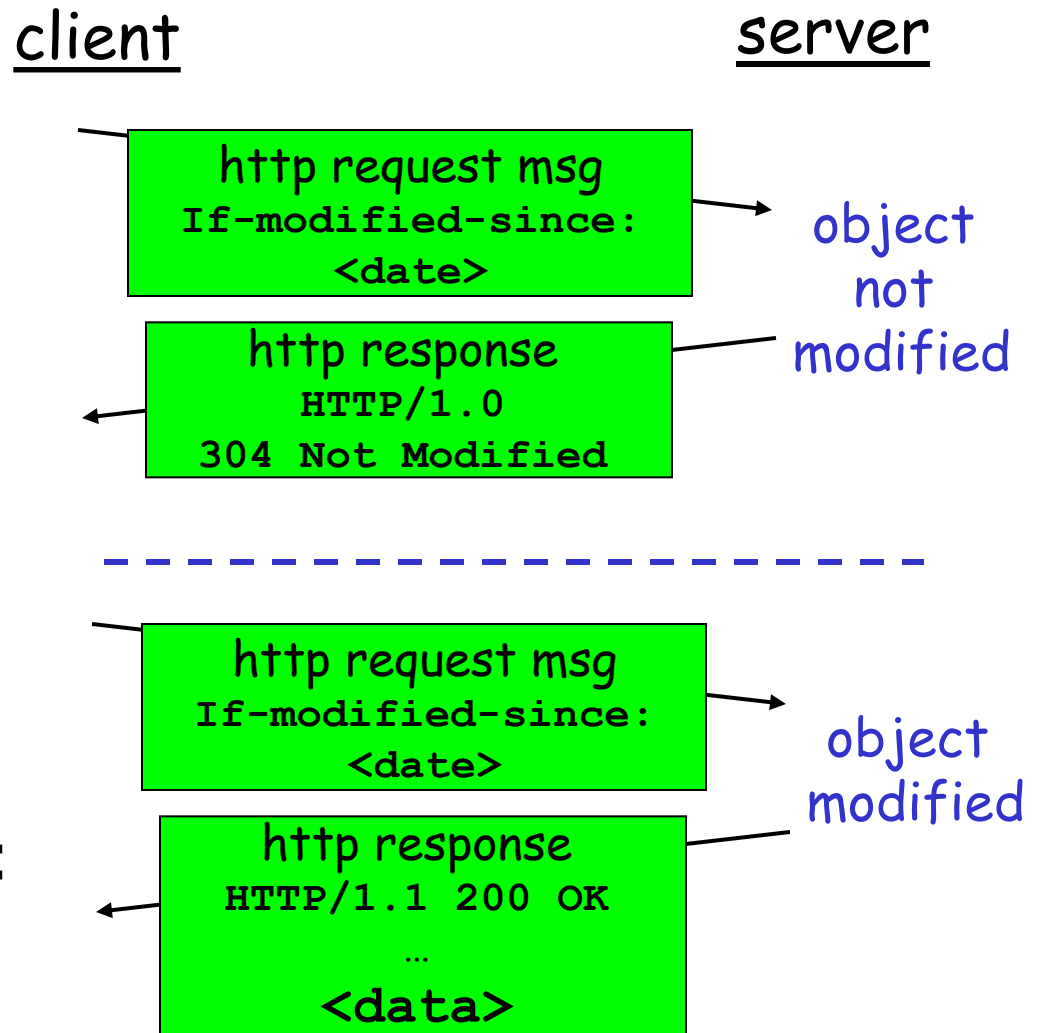


User-server interaction: Conditional GET

- ❑ **Goal:** Don't send object if client has up-to-date stored (cached) version
- ❑ Client: Specify date of cached copy in http request

`If-modified-since:`
`<date>`

- ❑ Server: Response contains no object if cached copy up-to-date:
`HTTP/1.0 304 Not Modified`



User-server state: Cookies

Most Web sites use cookies

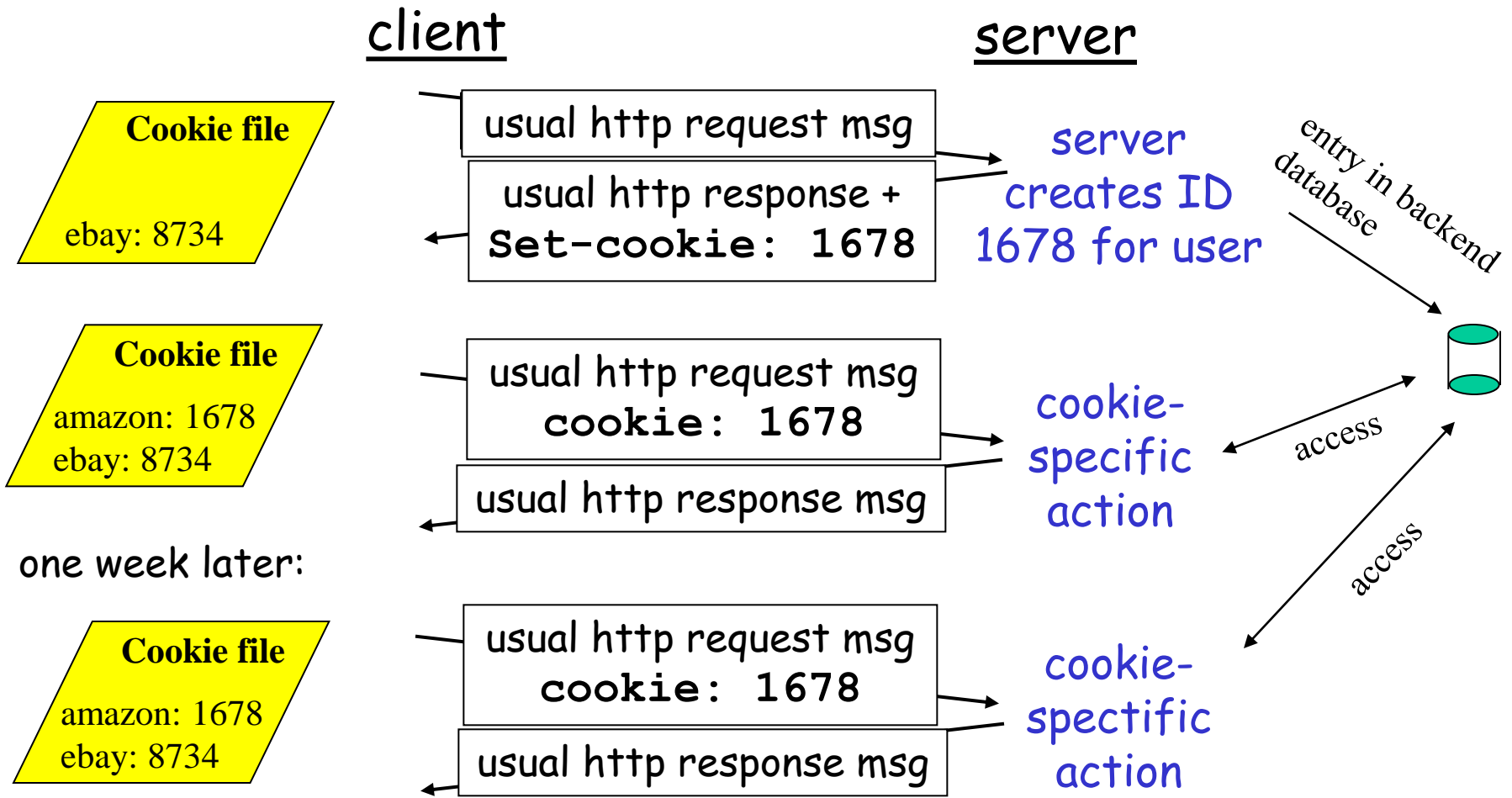
Four components:

- 1) Cookie header line of HTTP *response* message
- 2) Cookie header line in HTTP *request* message
- 3) Cookie file kept on user's host, managed by user's browser
- 4) Back-end database at Web site

Example:

- Susan access Internet always from same PC
- She visits a specific e-commerce site for first time
- When initial HTTP requests arrives at site, site creates a unique ID and creates an entry in backend database for ID

Cookies: Keeping "state"



Cookies: Keeping "state" (2)

What cookies can bring:

- ❑ Authorization
- ❑ Shopping carts
- ❑ Recommendations
- ❑ User session state
(Web e-mail)

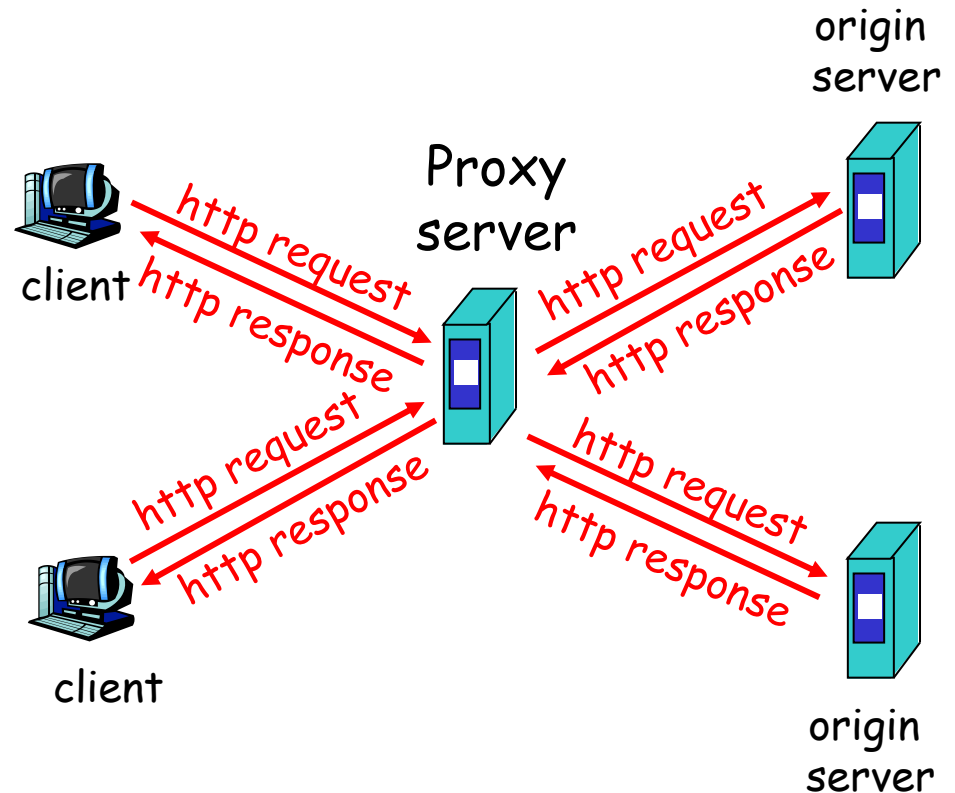
Cookies and privacy: aside

- ❑ Cookies permit sites to learn a lot about you
- ❑ You may supply name and e-mail to sites
- ❑ Search engines use redirection & cookies to learn yet more
- ❑ Advertising companies obtain info across sites

Web caches (proxy server)

Goal: satisfy client request without involving origin server

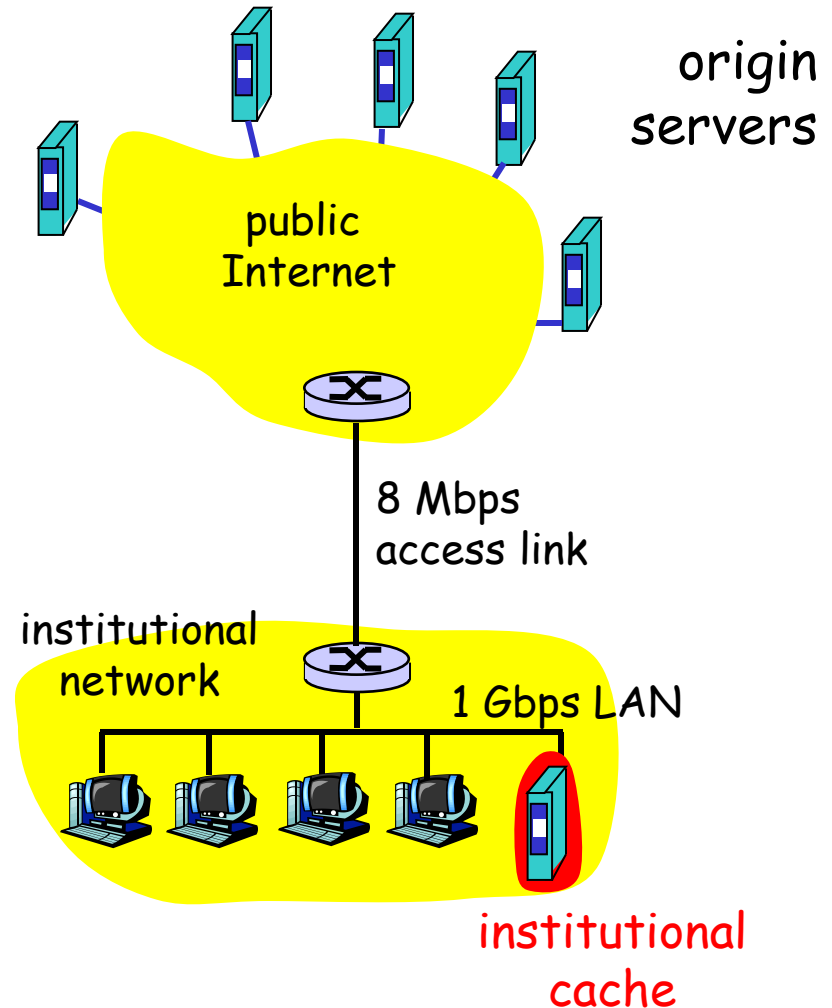
- User sets browser: WWW accesses via web cache
- Client sends all http requests to web cache
 - If object at web cache, web cache immediately returns object in http response
 - Else requests object from origin server, then returns http response to client



Why WWW caching?

Assume: Cache is “close” to client (e.g., in same network)

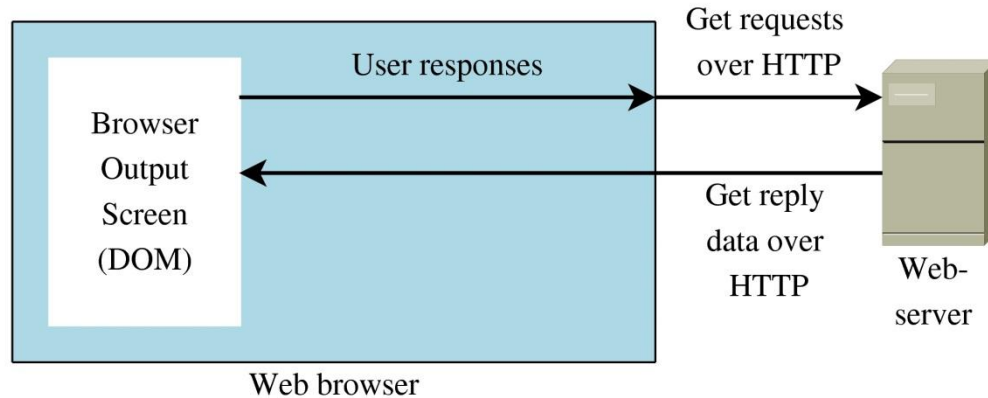
- ❑ Smaller response time: cache “closer” to client
- ❑ Decrease traffic to distant servers
 - Link out of institutional/local ISP network often bottleneck



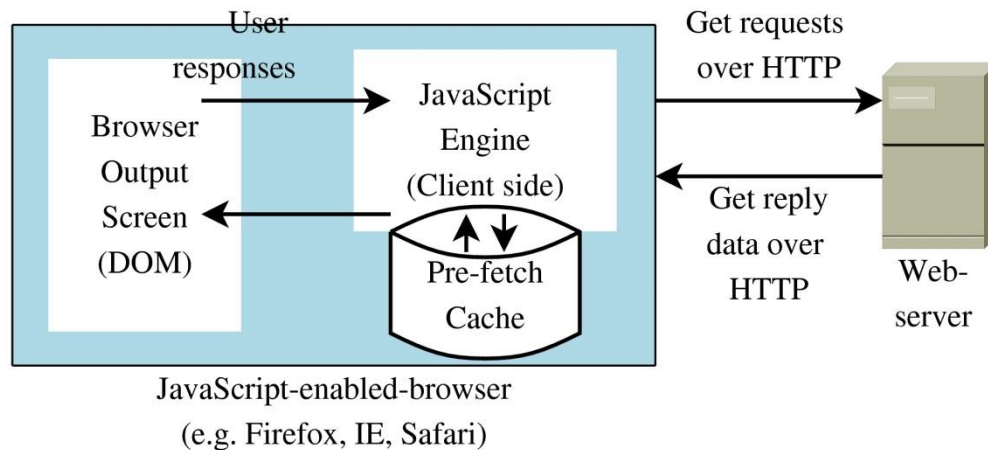
Web 2.0: E.g., AJAX enabled apps

- E.g.: Google Maps: a canonical AJAX application

(a) Classic
Web browsing



(b) AJAX enabled
Web browsing

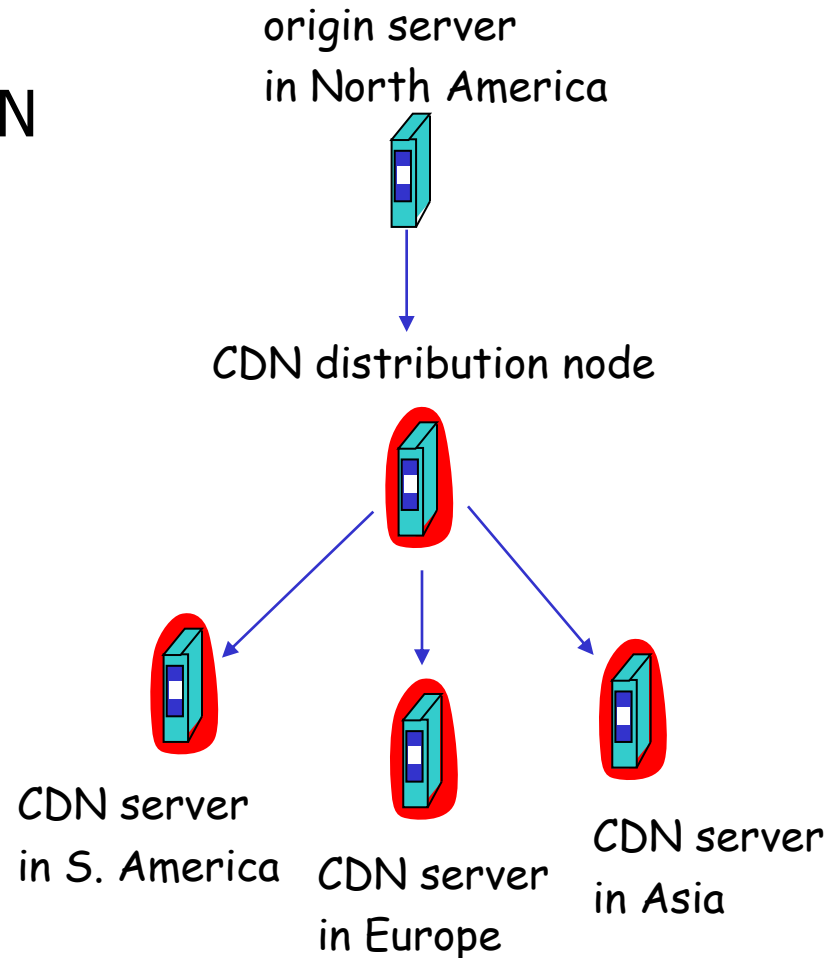


Content distribution networks (CDNs)

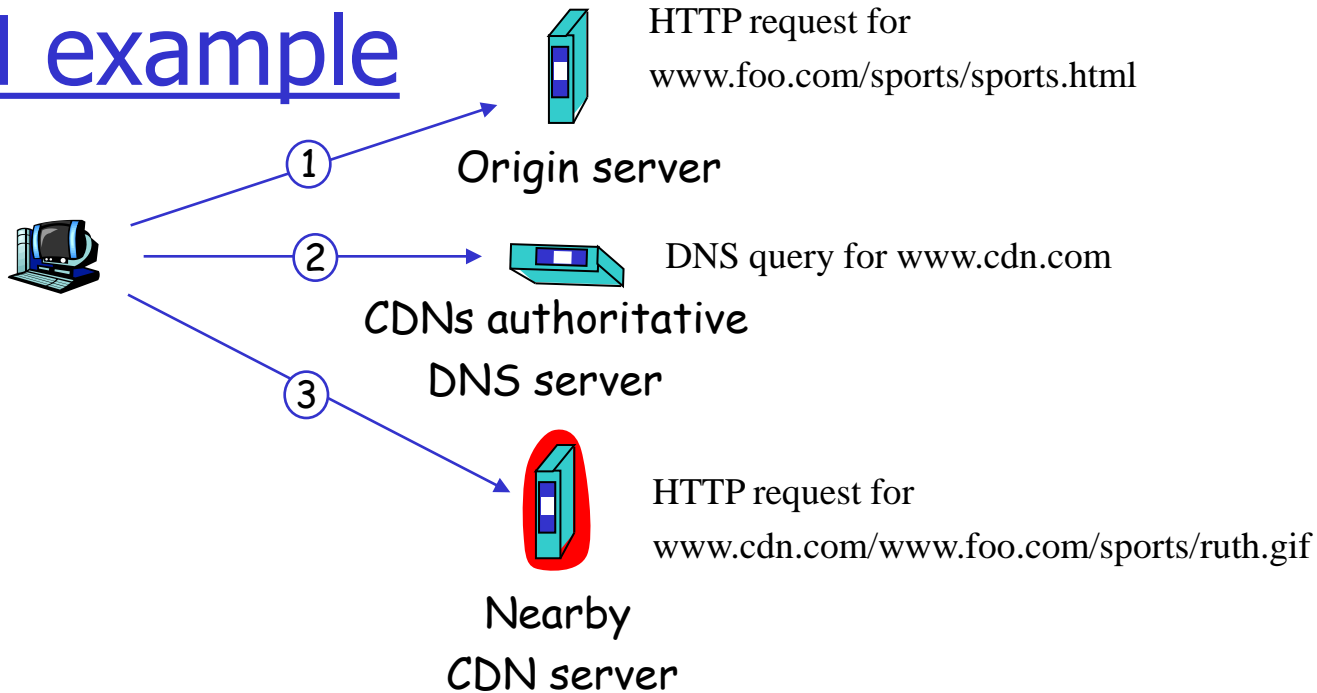
Content providers are the CDN customers.

Content replication

- ❑ CDN company installs hundreds of CDN servers throughout Internet
 - In lower-tier ISPs, close to users
- ❑ CDN replicates its customers' content in CDN servers. When provider updates content, CDN updates servers



CDN example



Origin server

- ❑ www.foo.com
- ❑ Distributes HTML
- ❑ Replaces:
http://www.foo.com/sports.ruth.gif
with
http://www.cdn.com/www.foo.com/sports/ruth.gif

CDN company

- ❑ cdn.com
- ❑ Distributes gif files
- ❑ Uses its authoritative DNS server to route redirect requests

More about CDNs

Routing requests

- ❑ CDN creates a “map”, indicating distances from leaf ISPs and CDN nodes
- ❑ When query arrives at authoritative DNS server
 - Server determines ISP from which query originates
 - Uses “map” to determine best CDN server

Not just Web pages

- ❑ Streaming stored audio/video
- ❑ Streaming real-time audio/video
 - CDN nodes create application-layer overlay network

DNS: Domain Name System

People: many identifiers:

- SSN, name, Passport #

Internet hosts, routers:

- IP address (32 bit) – used for addressing datagrams
- “name”, e.g., gaia.cs.umass.edu – used by humans

Q: Map between IP addresses and name?

□ Secure Domain Name System (DNS)

Dynamic Update: RFC 3007

DNS: Domain Name System

Domain Name System:

- ❑ *Distributed database:* Implemented in hierarchy of many *name servers*
- ❑ *Application-layer protocol:* Host, routers, name servers communicate to *resolve* names (address/name translation)
 - Core Internet function implemented as application-layer protocol
 - Complexity at network's "edge"

DNS name servers

Why not centralize DNS?

- ❑ Single point of failure
- ❑ Traffic volume
- ❑ Distant centralized database
- ❑ Maintenance

Does not *scale!*

DNS name servers (2)

No server has all name-to-IP address mappings

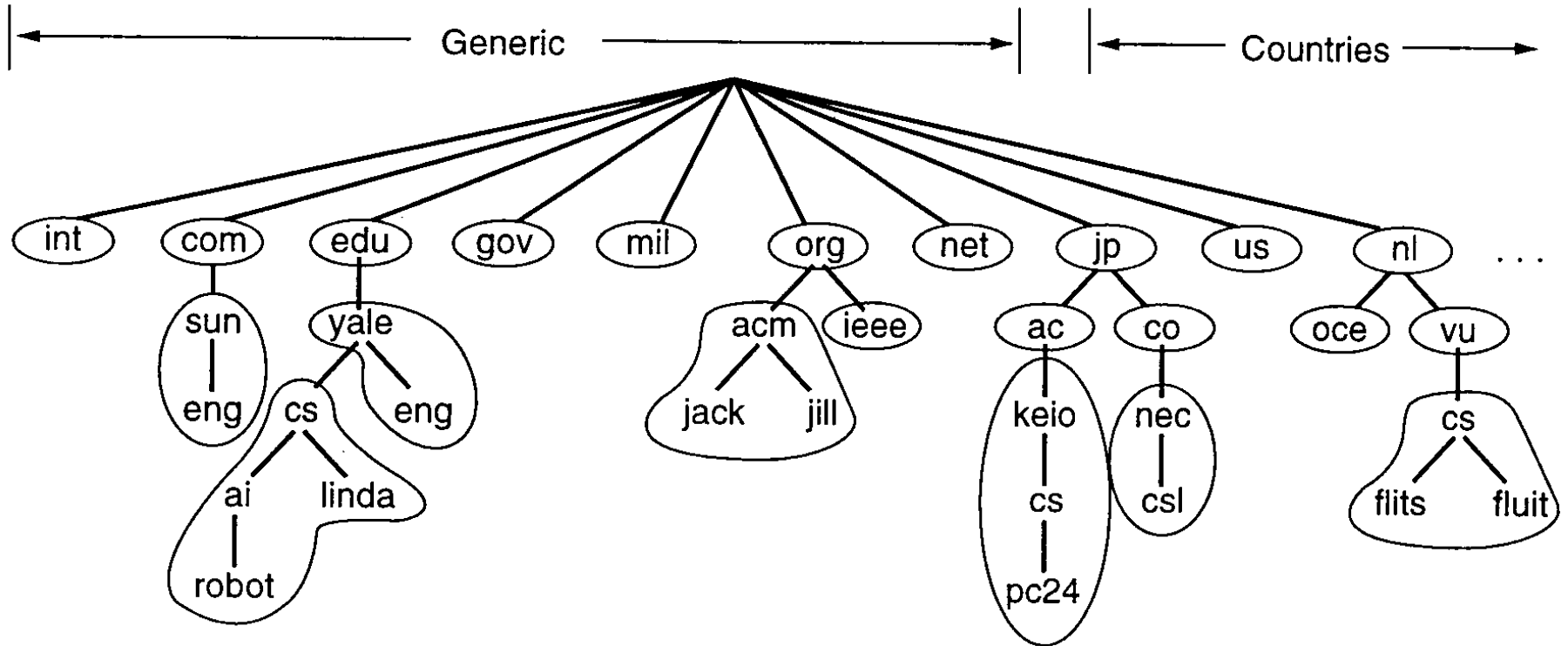
Local name servers:

- Each ISP, company has *local (default) name server*
- Host DNS query first goes to local name server

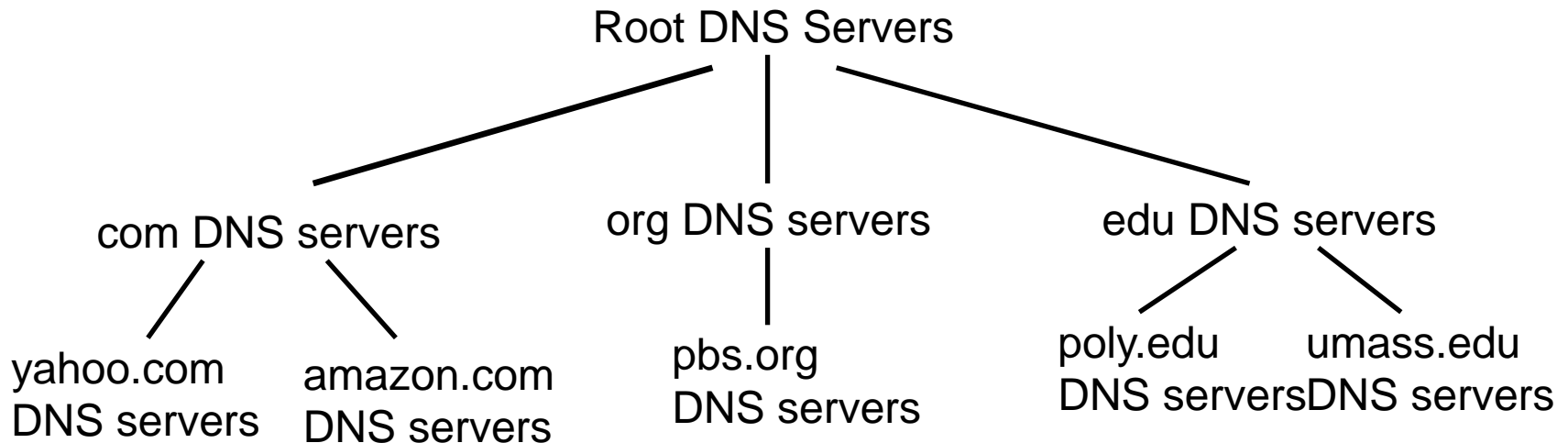
Authoritative name server:

- For a host: stores that host's IP address, name
- Can perform name/address translation for that host's name

DNS: Hierarchical naming tree



Distributed, hierarchical database

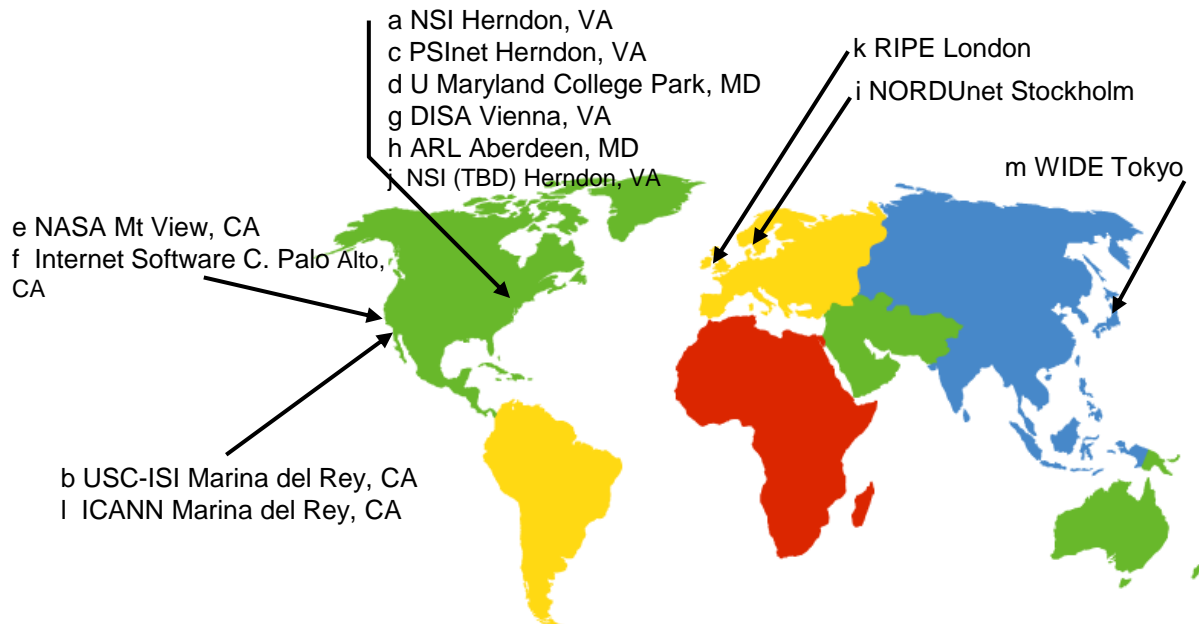


Client wants IP for www.amazon.com; 1st approx:

- ❑ Client queries a root server to find com DNS server
- ❑ Client queries com DNS server to get amazon.com DNS server
- ❑ Client queries amazon.com DNS server to get IP address for www.amazon.com

DNS: Root name servers

- ❑ Contacted by local name server that can not resolve name
- ❑ Root name server:
 - Contacts authoritative name server if name mapping not known
 - Gets mapping
 - Returns mapping to local name server
 - Some use anycast



3 root name servers worldwide

TLD and authoritative servers

- ❑ **Top-level domain (TLD) servers:** Responsible for com, org, net, edu, etc, and all top-level country domains uk, fr, ca, jp.
 - Network solutions maintains servers for com TLD
 - Educause for edu TLD
- ❑ **Authoritative DNS servers:** Organization's DNS servers, providing authoritative hostname to IP mappings for organization's servers (e.g., Web and mail).
 - Can be maintained by organization or service provider

Local name server

- ❑ Does not strictly belong to hierarchy
- ❑ Each ISP (residential ISP, company, university) has one.
 - Also called “default name server”
- ❑ When a host makes a DNS query, query is sent to its local DNS server
 - Acts as a proxy, forwards query into hierarchy.

DNS records

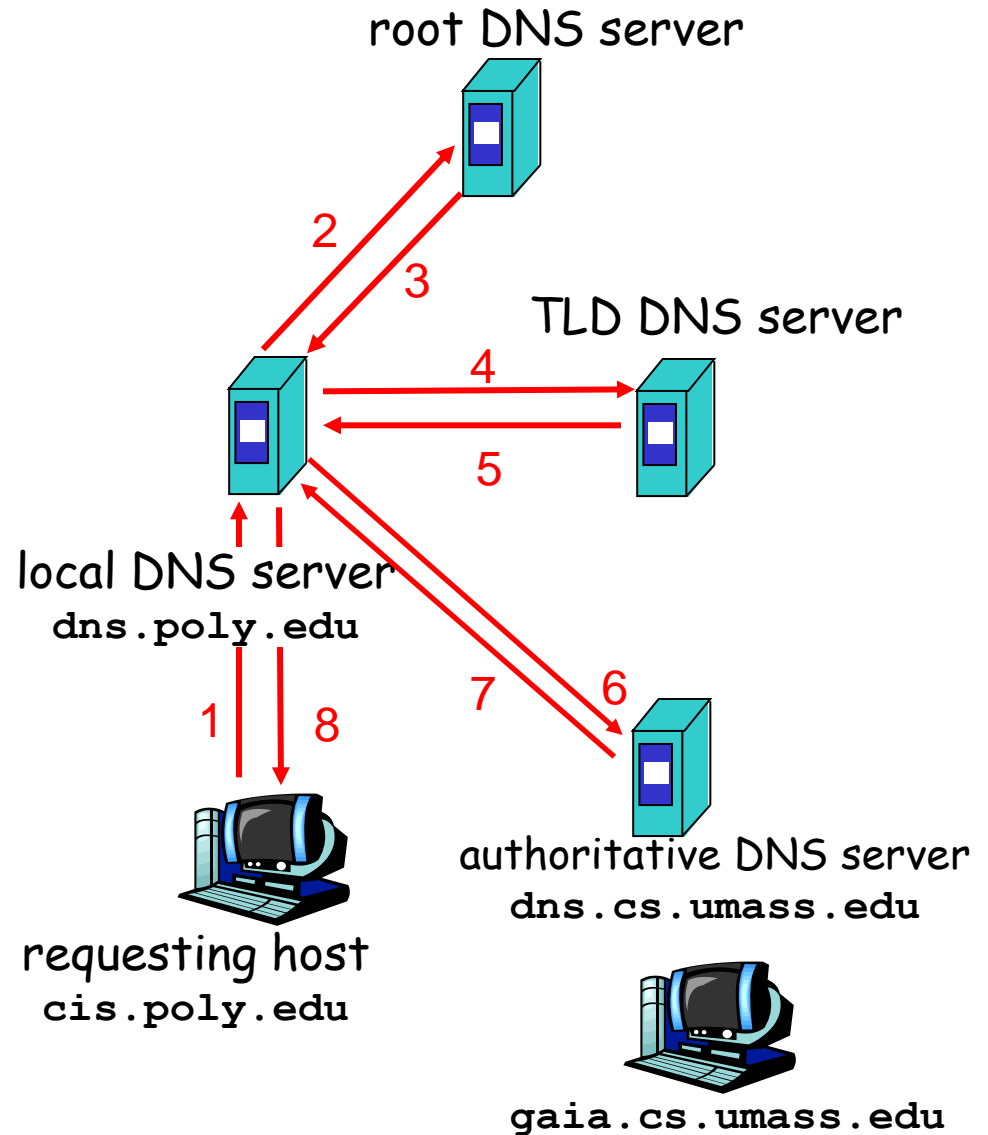
DNS: Distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

- ❑ Type=A
 - name is hostname
 - value is IP address
- ❑ Type=NS
 - name is domain (e.g., foo.com)
 - value is IP address of authoritative name server for this domain
- ❑ Type=CNAME
 - for alias
- ❑ Type=MX
 - for mail

Example

- Host at cis.poly.edu wants IP address for gaia.cs.umass.edu



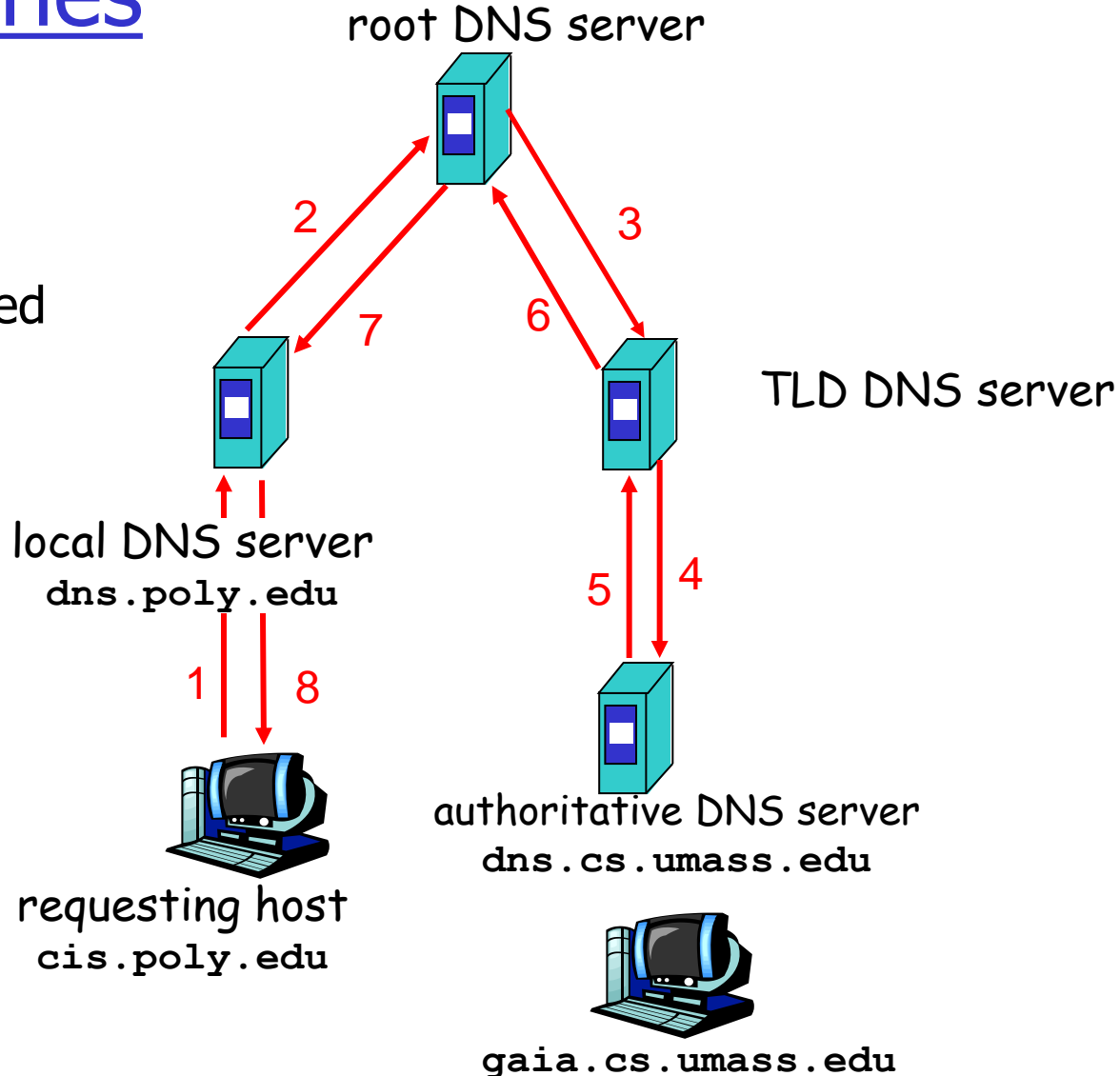
Recursive queries

Recursive query:

- ❑ Puts burden of name resolution on contacted name server
- ❑ Heavy load?

Iterated query:

- ❑ Contacted server replies with name of server to contact
- ❑ "I don't know this name, but ask this server"



DNS: Iterative queries

Recursive query:

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