

On the impact  
of variability on  
buffer dynamics  
in IP networks

# Mechanism that create variability

- ❑ User behavior
  - Application level variability, e.g., Web
- ❑ Network characteristics
  - Different delays, round-trip time, cross traffic
- ❑ Feedback control
  - Reliability and adaptivity

# Approach

- ❑ Study traces from simulations
  - Complete control over all aspects of network  
Workload model, Network model, Protocol
- ❑ Real network traces used as benchmark
  - Simulation setup and trace analysis
- ❑ Simulation traces used for trace driven simulation
  - Study queuing dynamic
- ❑ Correlation of analysis results to simulation setup
  - Identify network features that cause divergence

# Outline

- ❑ Simulation setup
  - Workload model
  - Network
  - Protocol
- ❑ Impact of variability at the application layer
  - Mice vs. elephants [V. Jacoson]
- ❑ Impact of feedback
  - Open loop vs. closed loop

# Workloads

## ❑ No variability:

- Infinite sources
  - 50 clients requesting big files

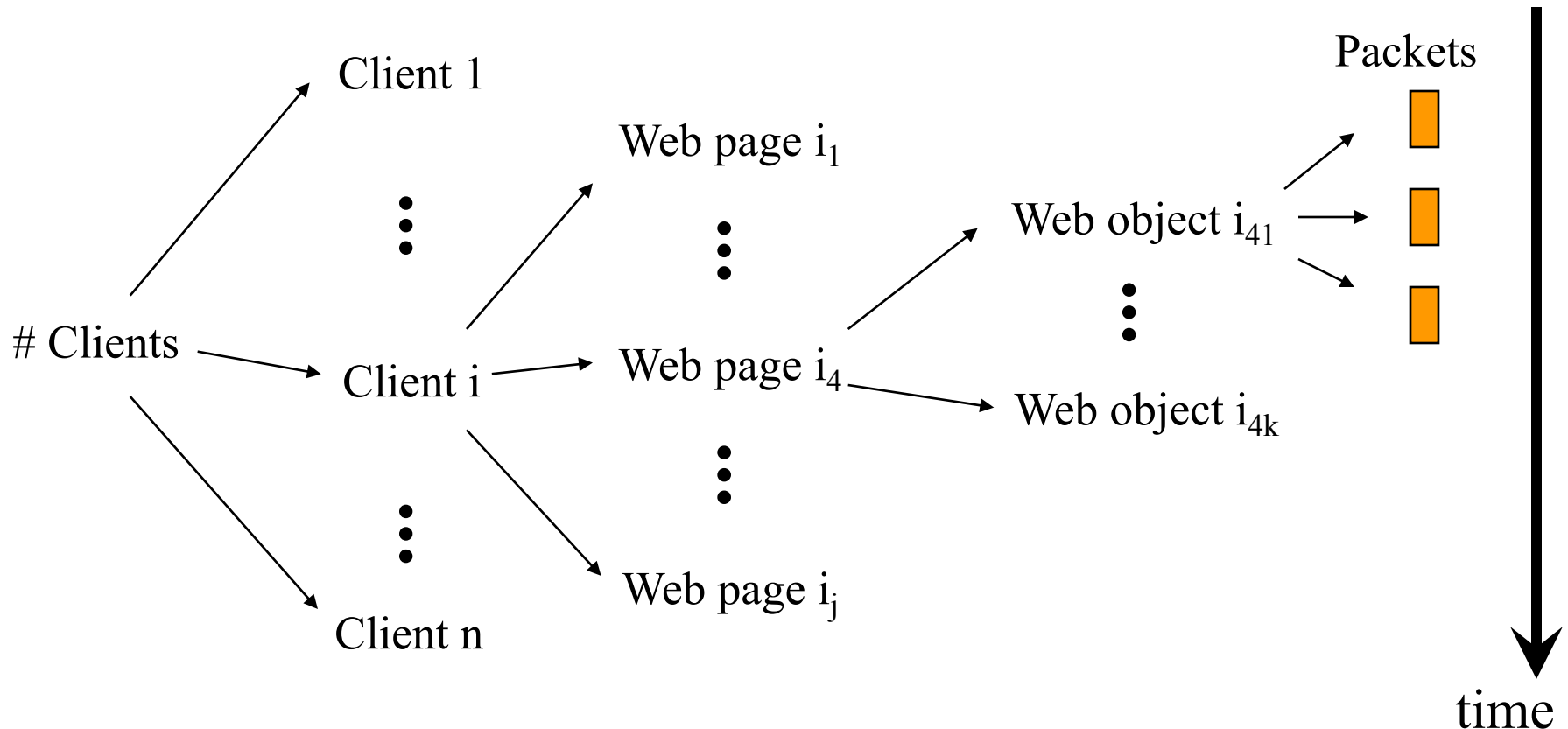
## ❑ High variability:

- Web sources
  - 350 clients down loading Web pages

## ❑ Simulation:

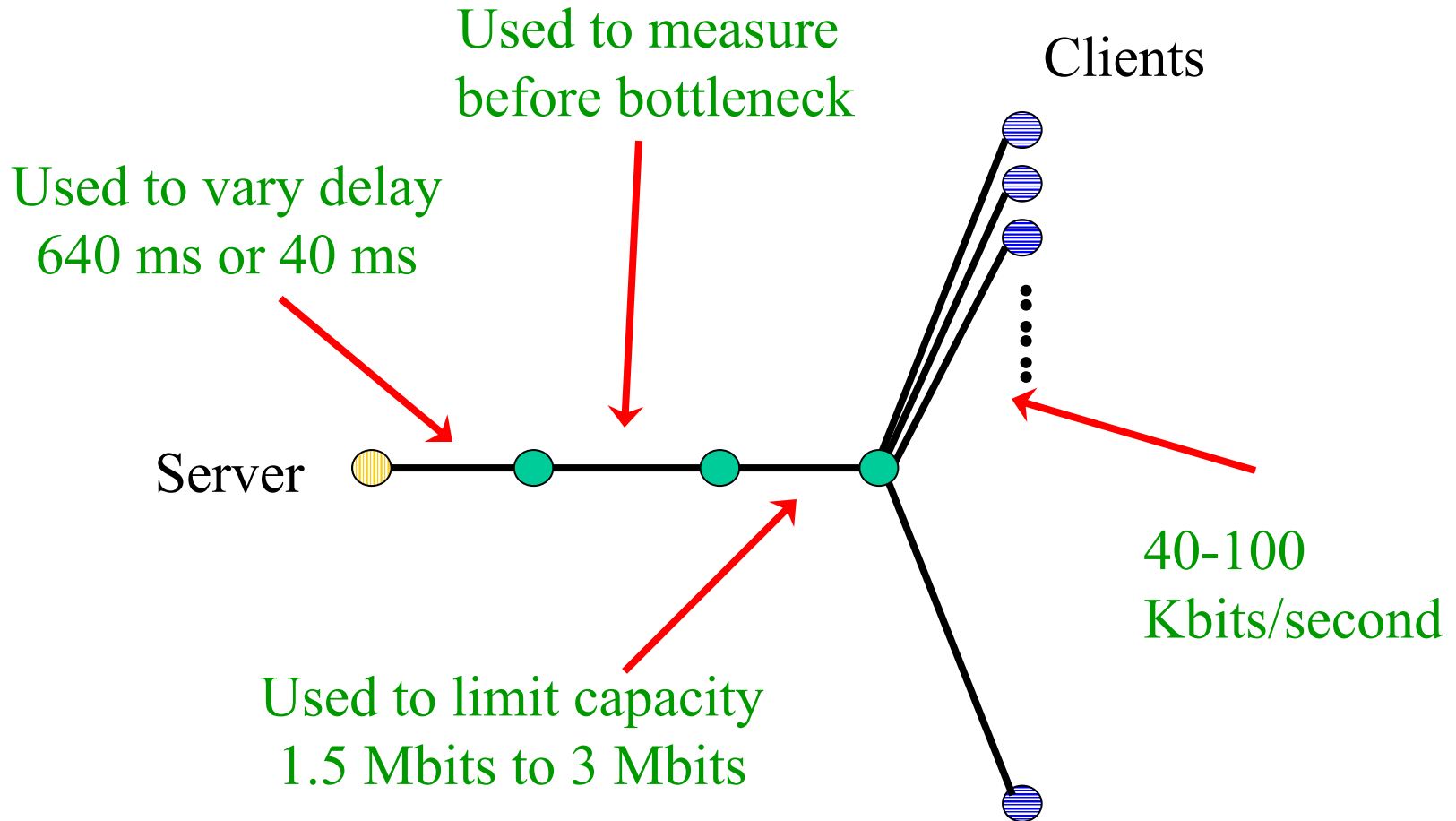
- Client startup: random time 0-600 seconds
- Duration: 4200 seconds
- Analysis: 900-4200 seconds

# High variability workload: Web



- Parameters (similar to SURGE [BC98]):
  - Number of clients, pages, objects, packets per object
  - Time between Web pages, Web objects

# A simple network topology



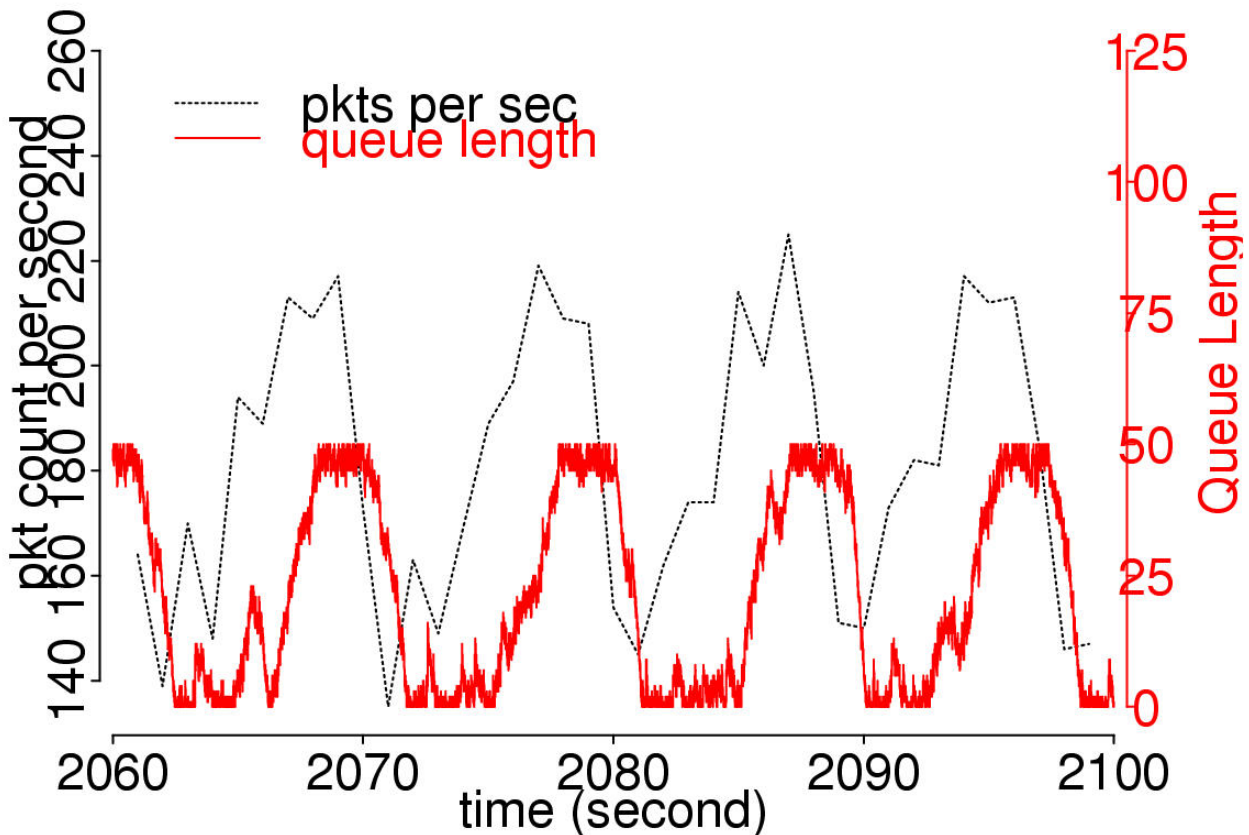
# TCP: Reliable byte stream

- ❑ Data segmented into segments
- ❑ Segments are acked by receiver (cumulative)
- ❑ Timer for every segment
- ❑ Segments retransmitted
  - Timer goes off
  - Four duplicate Acks received
- ❑ Flow control
  - Sliding window protocol avoids losses at receiver
  - Bandwidth limits impose congestion window
    - Slow start increases cwnd exponentially
    - Congestion avoidance increases cwnd linearly
    - Packet losses triggers congestion window changes



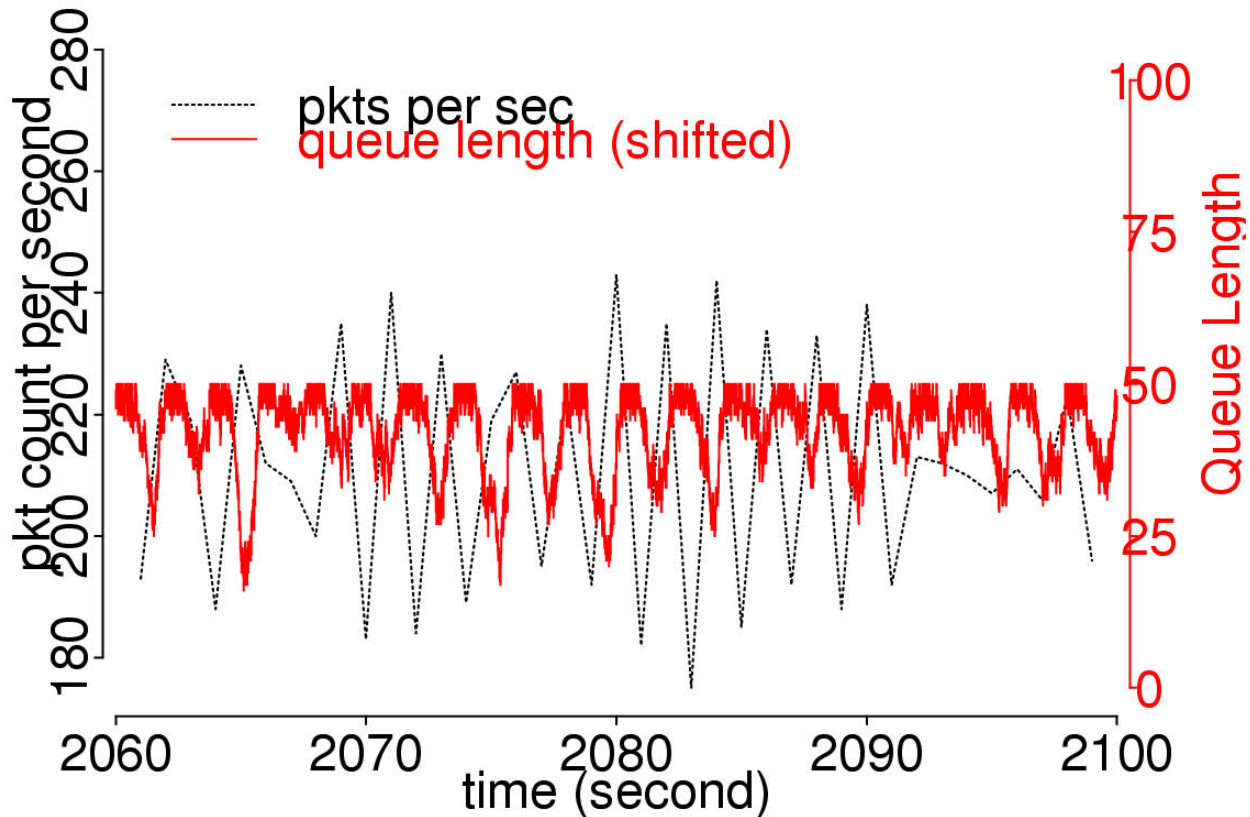
# Elephants: Infinite sources

- ❑ Packet rate process and buffer occupancy process
  - Network round trip time 1.3 seconds



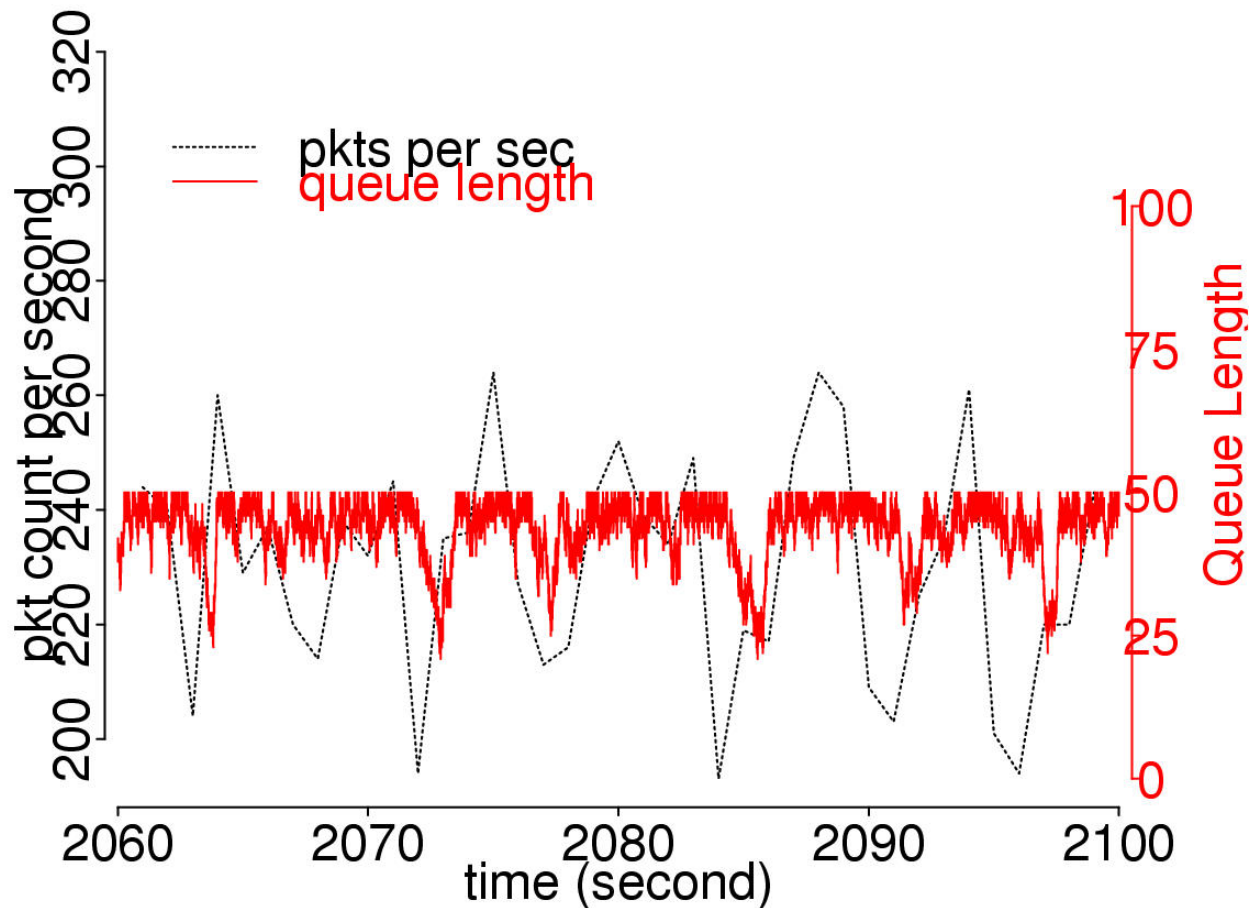
# Elephants (cont.)

- Packet rate process and buffer occupancy process
  - Network round trip time 0.14 seconds



# Mice and elephants: Web sources

- Significant portion of connections are short

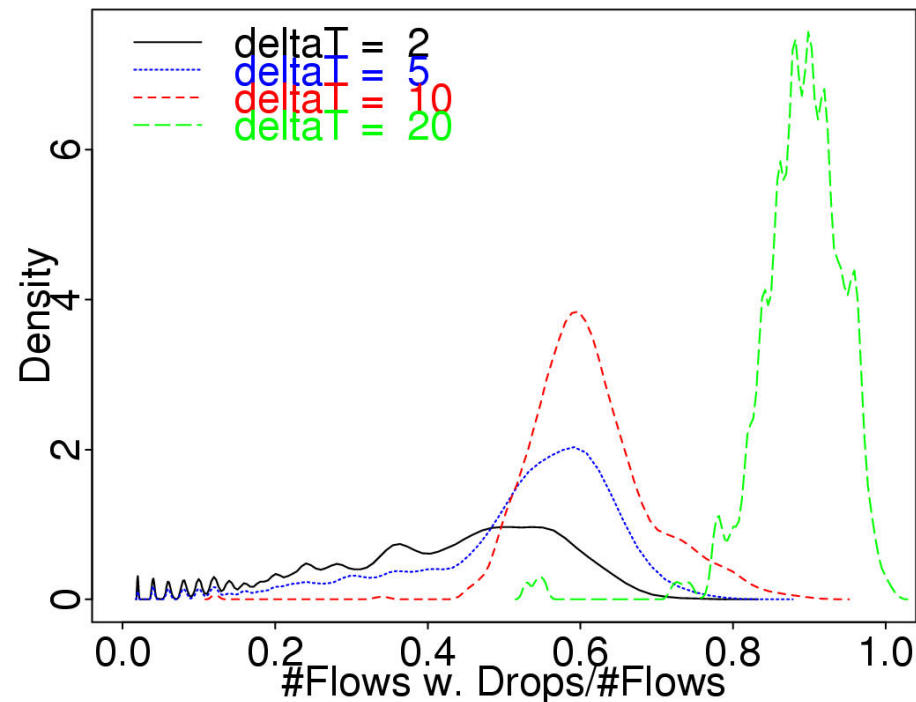


# Elephants vs. Mice

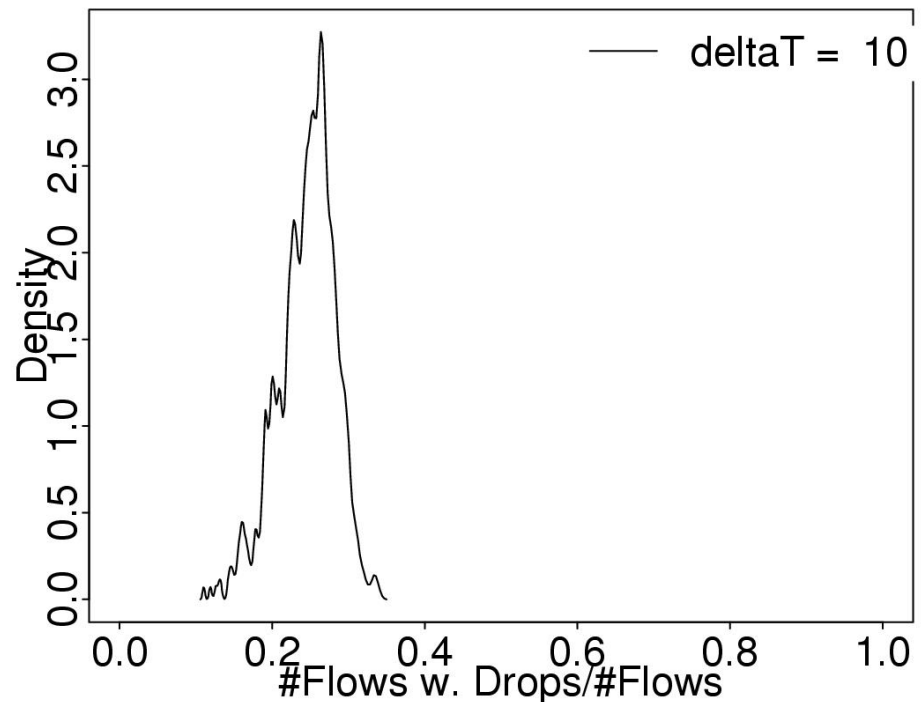
- ❑ No variability in workload and network
  - synchronization of packet rate process
  - synchronization of buffer occupancy process
- ❑ Mice stop synchronization
  - no apparent synchronization
  - higher packet arrival process
  - higher utilization

# Effect of synchronization

- Percentage of connections with losses during  $\Delta T$



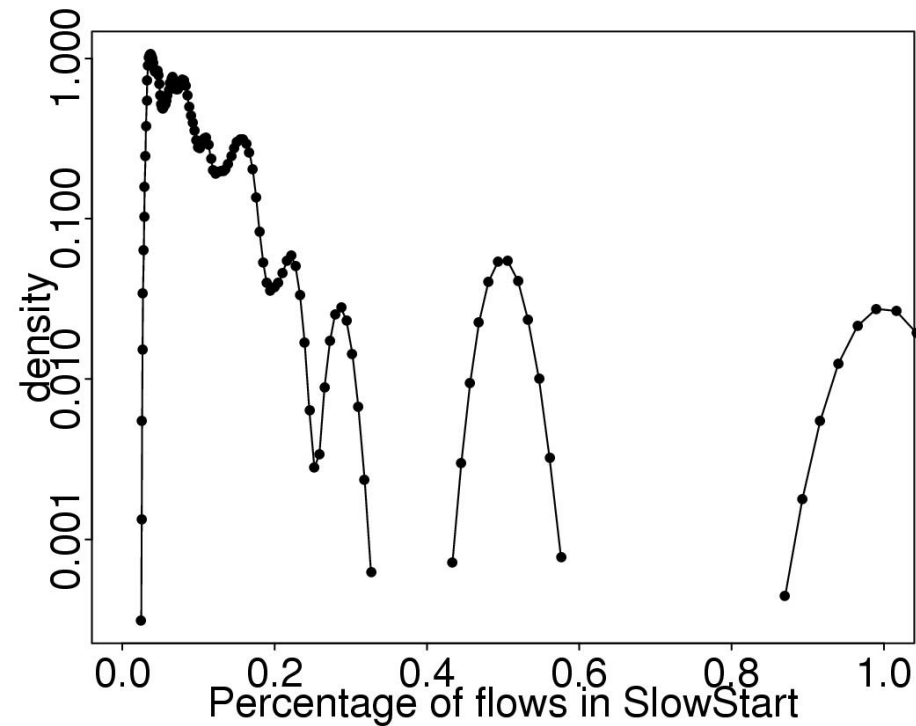
Infinite sources



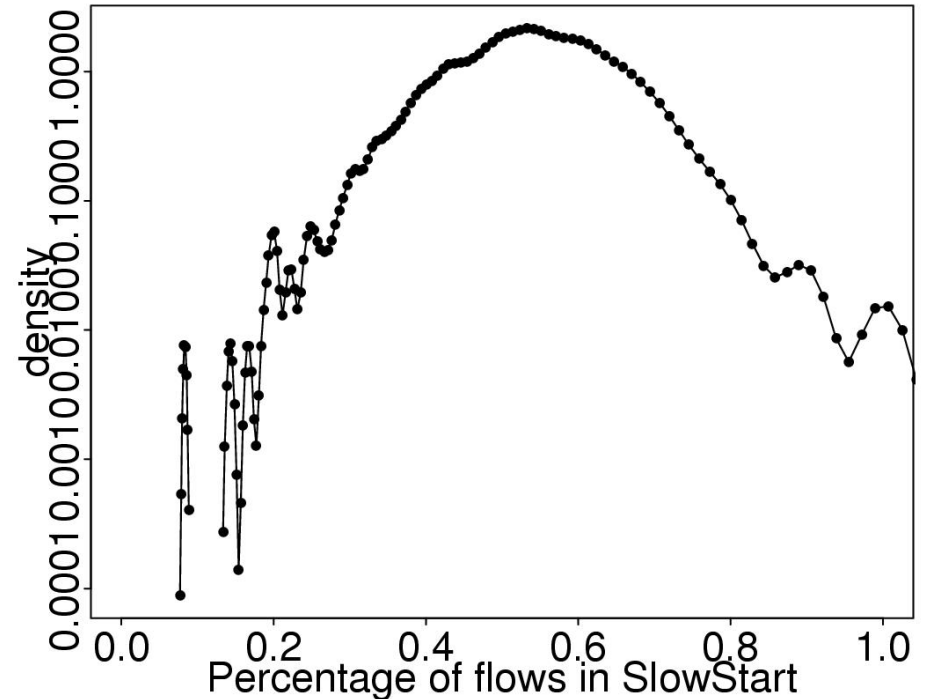
Web sources

# Effect of synchronization (cont.)

- Fractions of connections with losses in slow-start



Infinite sources



Web sources

# Why mice eliminate synchronization

## ❑ Mice

- Too short for feedback
- TCP states non-synchronized
- Arrival highly bursty
- Large fraction in slow-start

## ❑ Elephant

- Within two cycles losses for almost all connections
- TCP states synchronized
- Small percentage in slow start

## ❑ Consequence

- # of burst losses larger for Web than Infinite srcs

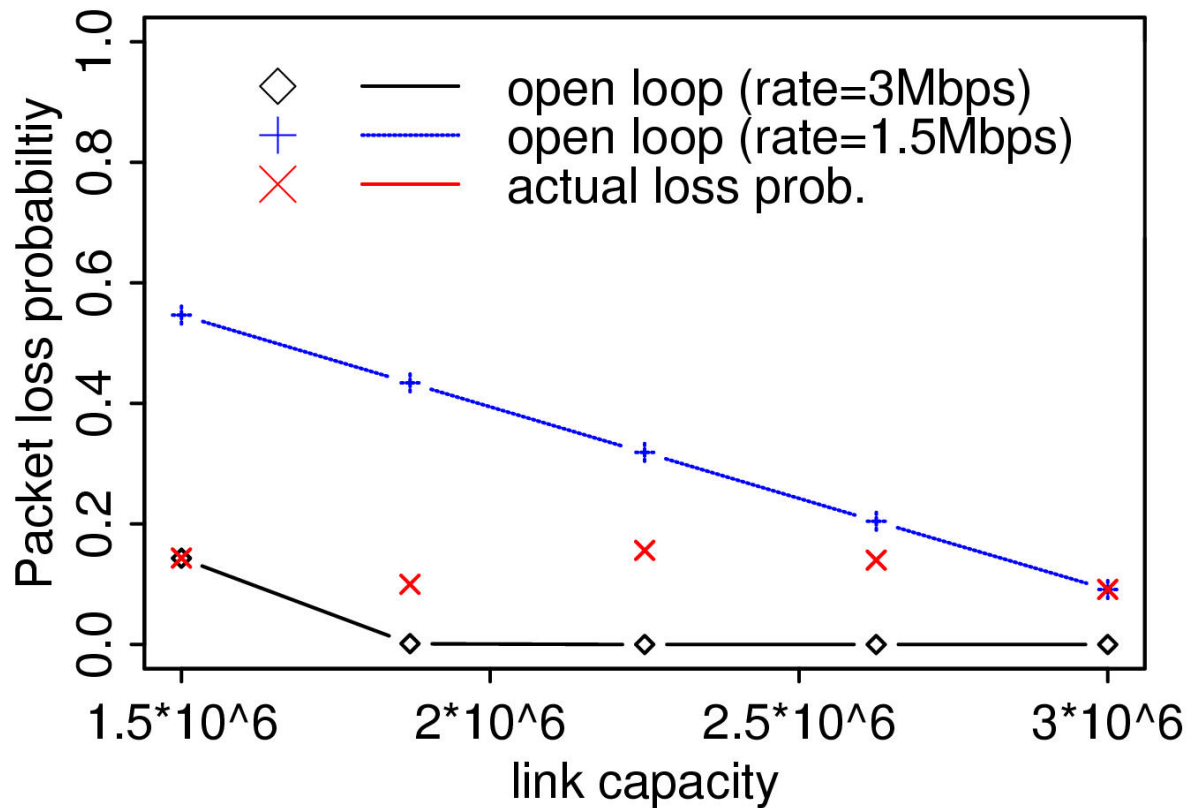
# Open loop vs. close loop

- ❑ Queuing system with finite buffer
  - Change in link capacity can be modeled by change of service time distribution
- ❑ A simple experiment
  - NS simulation with approx. finite buffer space (50 pkts)
  - Trace driven queuing analysis of G/D/1 queue with varying service times
- ❑ Evaluation
  - Set of ns-2 simulations with different bottleneck speeds
    - Calculate packet loss
  - Compare packet losses from simulations to packet loss predicted from open loop queue system



# Open loop vs. close loop (cont.)

## Web sources



- Either extremely conservative or overly aggressive

# Open loop assumptions

- ❑ Queuing system with infinite buffer
  - Buffer occupancy probability  $P[Q > x]$  can approximate finite buffer packet loss
- ❑ A simple experiment
  - ns-2 simulation (approx. inf. buffer space 1000 pkts)
  - Trace analysis to calculate buffer occupancy
- ❑ Evaluation
  - Set of ns-2 simulations with different finite buffers
  - Compare packet losses to buffer occupancy
- ❑ Result:
  - $P[Q > x]$  extremely conservative

# Conclusion

- ❑ Infinite source models and queue analysis provide necessary simplifications for
  - Analysis
  - Simulations
  
- ❑ Challenge
  - Address variability at
    - User level
    - Network level
  - Judge the impact of feedback